

Normal forms cont. & intro to transactions

superkey $X \rightarrow \{A_1, \dots, A_5\}$

A1	A2	A3	A4	A5

PK

name	location	salary
remy	LA	\$30
vincent	LA	\$20

first n.	last n.	location	salary
remy	w	LA	\$30
dan	s	seattle	\$50
dan	o	zurich	\$50

first n.	last n.	location	salary	course
remy	w	LA	\$30	143
remy	w	LA	\$30	240
remy	w	LA	\$30	249
dan	s	seattle	\$50	344
dan	s	seattle	\$50	444
dan	o	zurich	\$50	101
dan	o	zurich	\$50	113

name	job	location	salary	tax %
remy	prof	LA	\$30	20
dan	prof	seattle	\$50	15
vincent	TA	LA	\$20	10

$\text{job} \rightarrow \text{salary}$

$\text{name} \rightarrow \text{location} \implies \text{name, job} \rightarrow \text{tax \%}$

$\text{location, salary} \rightarrow \text{tax \%}$

BNCF

$$X \rightarrow Y : \begin{cases} Y \subseteq X \text{ (trivial FD)} \\ X \text{ is a superkey} \end{cases}$$

PK

name	location	salary
remy	LA	\$30
vincent	LA	\$20

first n.	last n.	location	salary
remy	w	LA	\$30
dan	s	seattle	\$50
dan	o	zurich	\$50

first n.	last n.	location	salary	course
remy	w	LA	\$30	143
remy	w	LA	\$30	240
remy	w	LA	\$30	249
dan	s	seattle	\$50	344
dan	s	seattle	\$50	444
dan	o	zurich	\$50	101
dan	o	zurich	\$50	113

name	job	location	salary	tax %
remy	prof	LA	\$30	20
dan	prof	seattle	\$50	15
vincent	TA	LA	\$20	10

job \rightarrow salary

name \rightarrow location \implies name, job \rightarrow tax %

location, salary \rightarrow tax %

Decomposition

Factor out violating FDs

$$X \rightarrow Y : \begin{cases} Y \subseteq X \text{ (trivial FD)} \\ X \text{ is a superkey} \end{cases}$$

Make new table over $X \cup Y$

Drop Y from old table (keep X)

first n.	last n.	location	salary	course
remy	w	LA	\$30	143
remy	w	LA	\$30	240
remy	w	LA	\$30	249
dan	s	seattle	\$50	344
dan	s	seattle	\$50	444
dan	o	zurich	\$50	101
dan	o	zurich	\$50	113

The diagram illustrates a table with five columns: **first n.**, **last n.**, **location**, **salary**, and **course**. The first two columns are highlighted with a solid orange border. A dashed orange border encloses the entire table. Two curved orange arrows originate from the top of the **first n.** and **last n.** columns and point to the **location** column, indicating a relationship between the first and last names and the location.

first n.	last n.	location	salary	course
remy	w	LA	\$30	143
remy	w	LA	\$30	240
remy	w	LA	\$30	249
dan	s	seattle	\$50	344
dan	s	seattle	\$50	444
dan	o	zurich	\$50	101
dan	o	zurich	\$50	113

first n.	last n.	location	salary
remy	w	LA	\$30
remy	w	LA	\$30
remy	w	LA	\$30
dan	s	seattle	\$50
dan	s	seattle	\$50
dan	o	zurich	\$50
dan	o	zurich	\$50

first n.	last n.	course
remy	w	143
remy	w	240
remy	w	249
dan	s	344
dan	s	444
dan	o	101
dan	o	113

job \rightarrow salary

name \rightarrow location

location, salary \rightarrow tax %

name	job	location	salary	tax %
remy	prof	LA	\$30	20
dan	prof	seattle	\$50	15
vincent	TA	LA	\$20	10

3rd Normal Form

$$X \rightarrow Y : \begin{cases} Y \subseteq X \text{ (trivial FD)} \\ X \text{ is a superkey} \\ \forall y \in Y : y \in \text{key} \end{cases}$$

3rd Normal Form

$$X \rightarrow Y : \begin{cases} Y \subseteq X \text{ (trivial FD)} \\ X \text{ is a superkey} \\ \forall y \in Y : y \in \text{key} \end{cases}$$

never loses FDs

decomposition in P time

3rd Normal Form

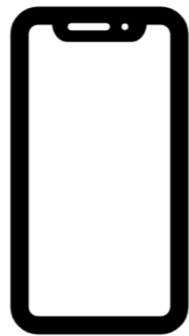
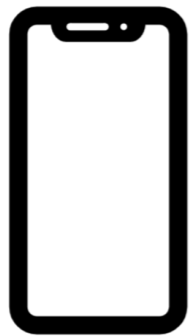
$$X \rightarrow Y : \begin{cases} Y \subseteq X \text{ (trivial FD)} \\ X \text{ is a superkey} \\ \forall y \in Y : y \in \text{key} \end{cases}$$

never loses FDs

decomposition in P time



Transactions: it's a wild world out there



exercise 1: free points for all!



atomicity

a TX either completes
or leaves no trace

exercise 2: playing favorites!



consistency

a TX should leave the DB

in a consistent state

exercise 3: armageddon!



isolation

multiple concurrent TX

should not interfere

exercise 4: blackout!

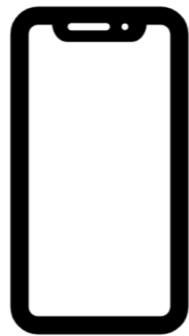
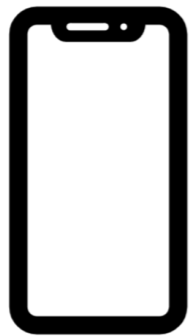


durability



completed TX are forever

SQLite demo



```
import sqlite3 as sql
```

```
con = sql.connect("bank.db")
```

```
cur = con.cursor()
```

```
res = cur.execute("SELECT * FROM acc")
```


transaction

a set of read & write ops

T1

READ(A, t)

t := t+100

WRITE(A, t)

READ(B, t)

t := t+100

WRITE(B,t)

A,B are
elements
in the DB

t is a local
variable
in the app

T1
READ (A, t) t := t+100 WRITE (A, t) READ (B, t) t := t+100 WRITE (B,t)

A,B are elements in the DB

t is a local variable in the app

T2
READ (A, s) s := s*2 WRITE (A,s) READ (B,s) s := s*2 WRITE (B,s)

transaction

a set of read & write ops

A either execute all, or nothing

C

I 2 transactions do not interfere

D

schedule

an interleaving & read/writes

from different TXs

time



T1	T2
READ(A, t)	READ(A, s)
	$s := s^2$
$t := t+100$	
WRITE(A, t)	WRITE(A, s)
	READ(B, s)
	$s := s^2$
READ(B, t)	
	WRITE(B, s)
$t := t+100$	
WRITE(B, t)	

serial schedule

"one at a time!"

time



A = 2
B = 2

T1

READ(A, t)

t := t+100

WRITE(A, t)

READ(B, t)

t := t+100

WRITE(B,t)

T2

READ(A, s)

s := s*2

WRITE(A,s)

READ(B,s)

s := s*2

WRITE(B,s)

time



A = 2
B = 2

A = 4
B = 4

A = 104
B = 104

T1

T2

READ(A, t)
t := t+100
WRITE(A, t)
READ(B, t)
t := t+100
WRITE(B, t)

READ(A, s)
s := s*2
WRITE(A, s)
READ(B, s)
s := s*2
WRITE(B, s)

serializable schedule

equivalent to "one at a time!"

T1	T2	
READ (A, t)		A = 2 B = 2
t := t+100		
WRITE (A, t)		
	READ (A, s)	A = 102 B = 2
	s := s*2	
	WRITE (A,s)	A = 204 B = 2
READ (B, t)		
t := t+100		
WRITE (B,t)		
	READ (B,s)	A = 204 B = 102
	s := s*2	
	WRITE (B,s)	A = 204 B = 204

serial → isolation

serializable = serial

serializable → isolation