

# Isolation Levels & Application Programming

# isolation levels

**FAST**

**CORRECT**



None (Chaos?)

Read Uncommitted

Read Committed

Repeatable Read

Serializable

# isolation levels

avoid certain types of problems:

dirty read/inconsistent read

lost update

unrepeatable read

# **dirty/inconsistent read**

seeing updates from uncommitted TXN

*Manager wants to  
balance project budgets*

*CEO wants to check  
company balance*

SELECT SUM(budget) ...

-\$10mil from project A



+\$7mil to project B



+\$3mil to project C



time



# **lost update**

update overwritten by another TXN

$$A = B = 100$$

$W(A, 200)$	
$W(B, 0)$	
	$W(B, 200)$
	$W(A, 0)$

$$A + B = 200 \checkmark$$

$$A = B = 100$$

$W(A, 200)$	
	$W(B, 200)$
$W(B, 0)$	
	$W(A, 0)$

$$A = B = 0 \text{ X}$$



# **unrepeatable read**

two reads give different results

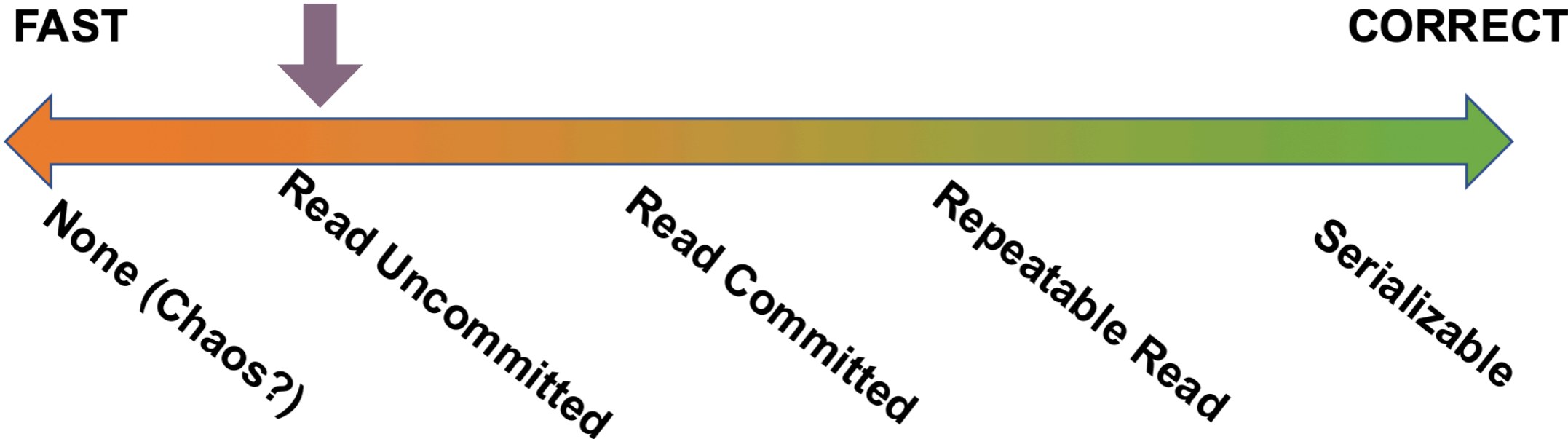
$A = 100$

	$R(A)$
$W(A, 0)$	
	$R(A)$

$A = 100$

$A = 0$

# isolation levels



# **read uncommitted**

Strict 2PL for writes

no locks at all for reads!

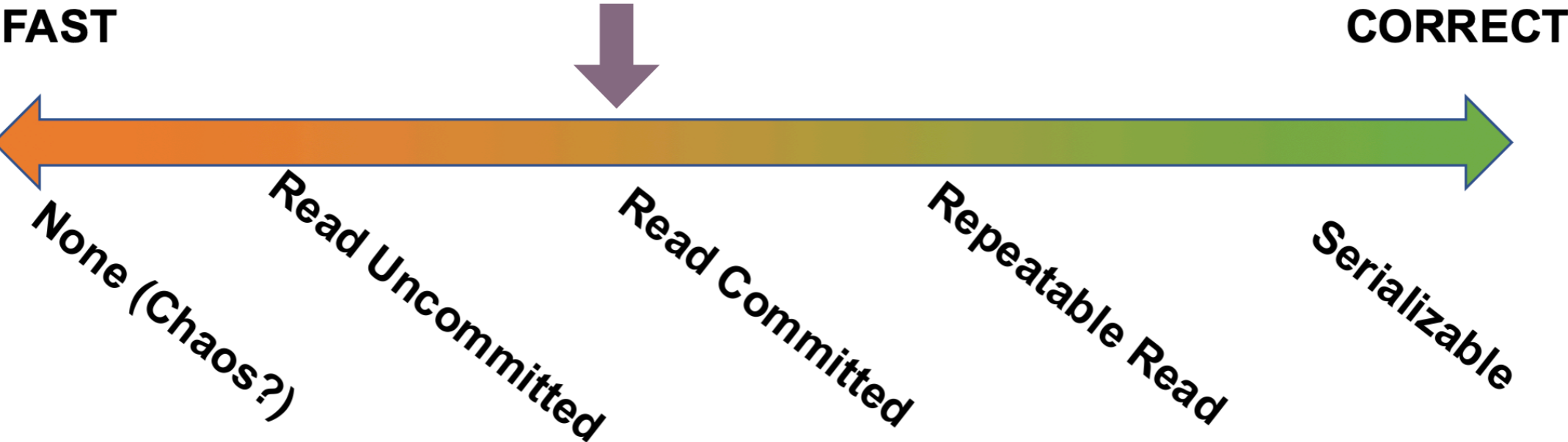
# **read uncommitted**

very fast reads

assumes few/no writes

read accuracy is not critical

# isolation levels



# read committed

Strict 2PL for writes

on-demand read locks (not 2PL!)

lock  $\rightarrow$  R  $\rightarrow$  unlock

**no dirty reads,** possible unrepeatable reads

## no dirty reads

$W(A, 0)$	
	<del><math>L(A), R(A)</math></del>
COMMIT, $U(A)$	



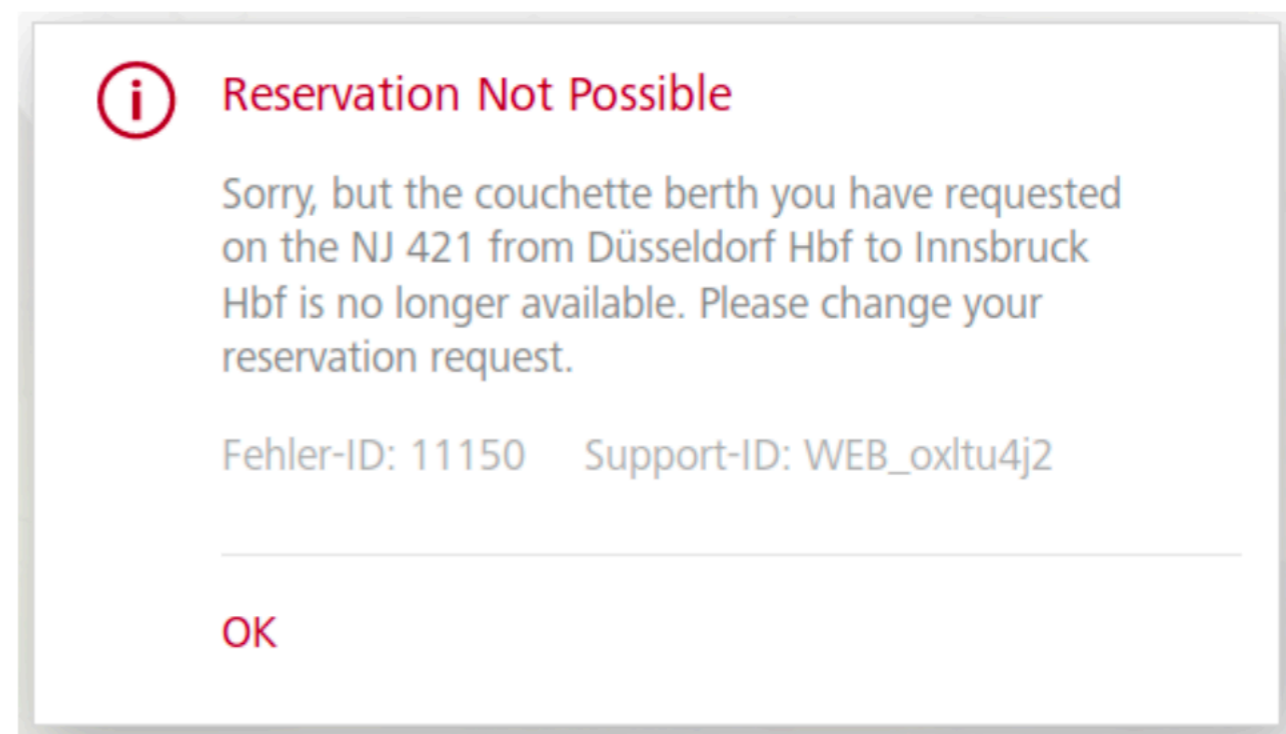
# possible unrepeatable reads

	L(A),R(A),U(A)
L(A),W(A),U(A)	
	L(A),R(A),U(A)

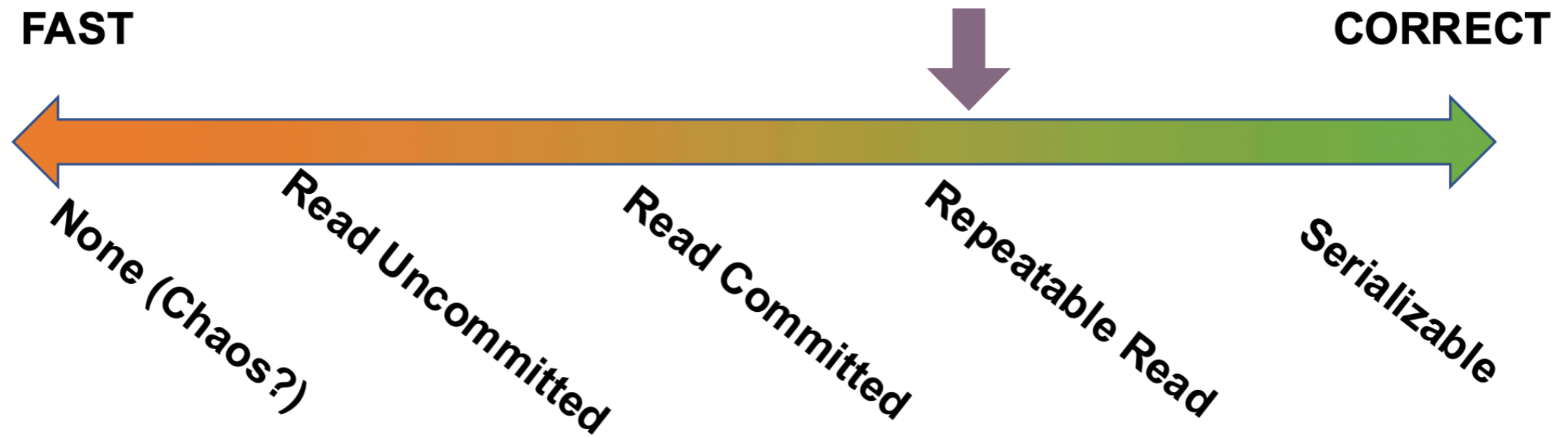
# read committed

guarantee read result is valid at *some* point

useful for online shops



# isolation levels



# **repeatable read**

Strict 2PL write locks

Strict 2PL read locks

conflict serializable!

but not serializable???

# The Phantom Menace

- Same read has more rows
- Asset checking scenario:

Accountant wants to check company assets

```
SELECT *  
FROM products  
WHERE price < 10.00
```

```
SELECT *  
FROM products  
WHERE price < 10.00
```

Warehouse catalogs new products

```
INSERT INTO Products  
VALUES ('nuts', 10, 8.99)
```

time



# the phantom problem

SELECT *	R(A), R(B)	
		W(C)
SELECT *	R(A), R(B), R(C)	

INSERT

# the phantom problem

conflict serializable → serializable **w/o inserts**

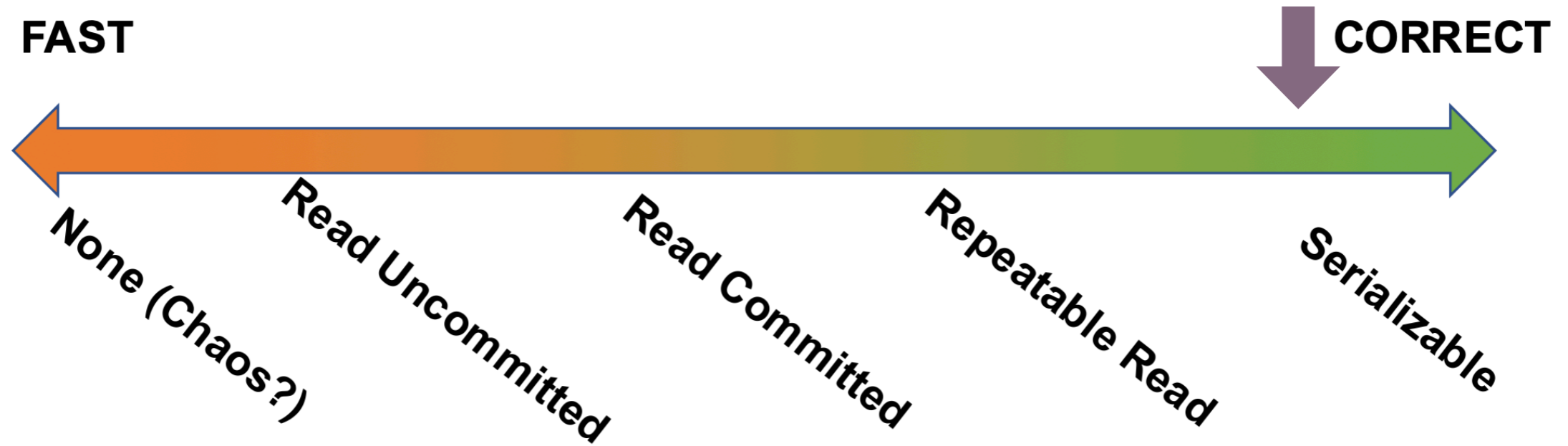
solution: lock entire table

# the phantom problem

SELECT *	<b>L(T)</b> , R(T)	
		<del>L(T), W(C)</del>
SELECT *	R(T), C, U(T)	



# isolation levels



HI, THIS IS  
YOUR SON'S SCHOOL.  
WE'RE HAVING SOME  
COMPUTER TROUBLE.



OH, DEAR - DID HE  
BREAK SOMETHING?

IN A WAY - )



DID YOU REALLY  
NAME YOUR SON  
Robert'); DROP  
TABLE Students;-- ?



OH, YES. LITTLE  
BOBBY TABLES,  
WE CALL HIM.

WELL, WE'VE LOST THIS  
YEAR'S STUDENT RECORDS.  
I HOPE YOU'RE HAPPY.



AND I HOPE  
YOU'VE LEARNED  
TO SANITIZE YOUR  
DATABASE INPUTS.

```
INSERT INTO students (id, 'name');
```

id	name
111	...
123	Robert'); DROP TABLE Students;--
145	...

```
INSERT INTO students (id, 'Robert');  
DROP TABLE Students;--');
```

```
SELECT * FROM user
WHERE name = 'x'
AND password = 'y'
```

**attack:** return all users

**hint:** use UNION / OR

# **solution to SQL injection**

parameterized queries

prepared statements

access control



# **passwords**

NEVER store passwords in plain text!

# hash functions

deterministic

$$x = y \implies h(x) = h(y)$$

low collision

$$x \neq y \implies h(x) \neq h(y)$$

$$(P(h(x) = h(y)) \approx 0)$$

easy to compute

$$f(x) \in O(1)$$

hard to invert

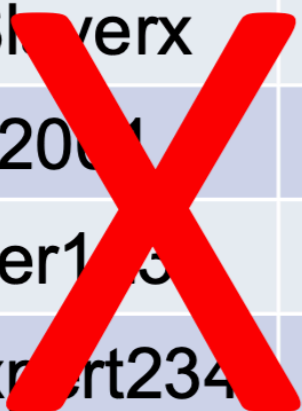
$$f^{-1}(x) \in O(2^N)$$

# **passwords**

NEVER store passwords in plain text!

**store hash instead**

Username	Password
bobtheninja246	password
xDragonSlayerx	asdf
annabelle2001	password
lamamaster123	ilovefish
theSQLexpert234	j62ld12446
seahawksrule12	j62ld12446

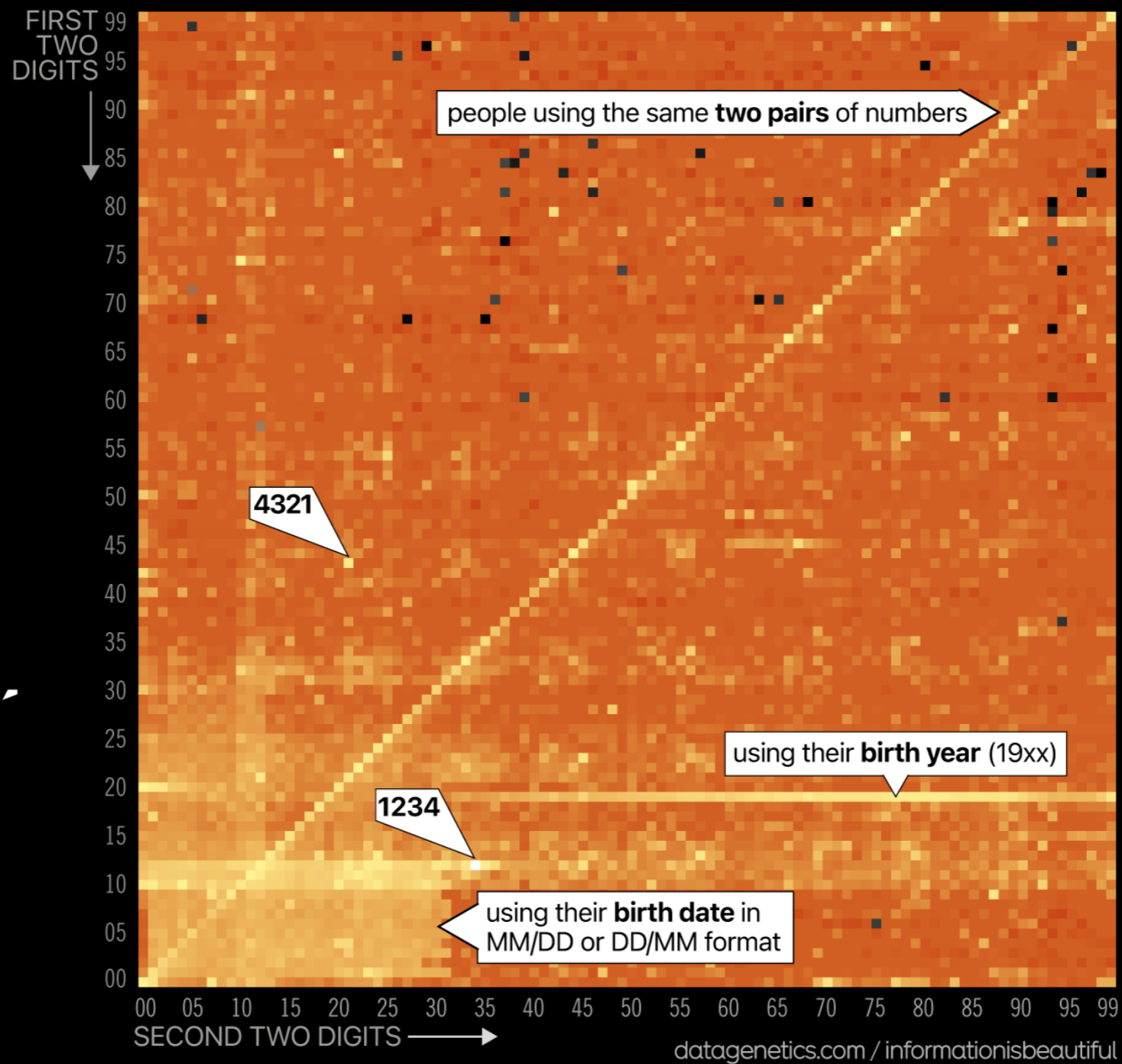


Username	HashedPassword
bobtheninja246	3da541...
xDragonSlayerx	bfd361...
annabelle2001	3da541...
lamamaster123	5baa61...
theSQLexpert234	ca8612...
seahawksrule12	ca8612...

```
SELECT * FROM user  
WHERE name = 'x'  
AND pw_hash = hash('y')
```

# Most to Least Common 4-Digit PIN Numbers

sourced from multiple data breaches



## most common

1234 0000 7777 2000 2222 9999 5555 1122 8888 2001  
 1111 1212 1004 4444 6969 3333 6666 1313 4321 1010

27% of all PIN numbers

## least common

8557 8438 9539 7063 6827 0859 6793 0738 6835 8093  
 9047 0439 8196 6093 7394 9480 8398 7637 9629 8068

Username	Password
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# **people are bad at passwords!**

reuse pw across sites

commonly used pw



```
salt = getRandom()  
salted_pw_hash = hash(pw, salt)
```

Username	Password
bobtheninja246	<b>password</b>
xDragonSlayerx	<b>asdf</b>
annabelle2001	password
lamamaster123	ilovefish
theSQLexpert234	j62ld12446
seahawksrule12	j62ld12446



Username	Salt	HashedPassword
bobtheninja246	17	<b>7a4959...</b>
xDragonSlayerx	m9	<b>59438a...</b>
annabelle2001	23	4c812e...
lamamaster123	q7	3e0e04...
theSQLexpert234	k3	dcfea6...
seahawksrule12	ji	e840fc...

# **privacy laws**

HIPPA

GDPR

FERPA

## **common approach**

"de-ID": remove personal identifiable information

## common approach

"de-ID": remove personal identifiable information



Latanya Sweeney: **not enough!**

Re-identification of Mass. Gov. William Weld

```
health_rec(zip, DoB, sex,  
           diagnosis, procedure, ...)
```

```
voter(name, party, ...,  
      zip, DoB, sex, ...)
```

```
("Weld", "R.", ...,  
 12345, 02-30, M, ...)
```



Cambridge, MA Voter Data (\$20)

Name	ZIP	Sex	Bday
...	...	...	...
W. Weld	12345	M	Feb 30
...	...	...	...



Anon. Insurance Data for Researchers

ZIP	Sex	Bday	MedInfo
...	...	...	...
12345	M	Feb 30	Afluenza
...	...	...	...

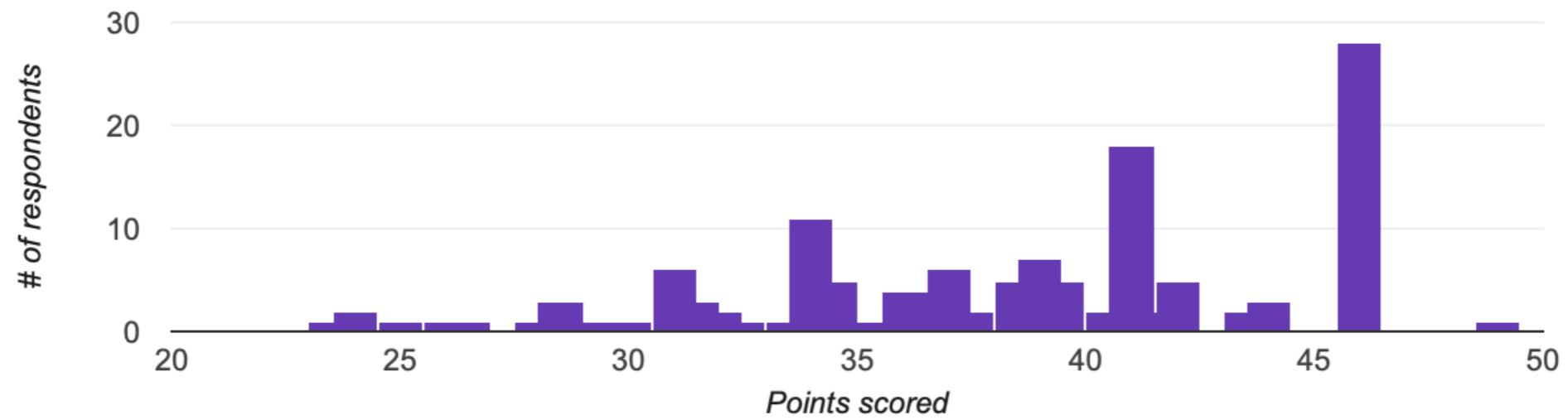


**Average**  
38.45 / 49 points

**Median**  
39 / 49 points

**Range**  
23.5 - 49 points

**Total points distribution**

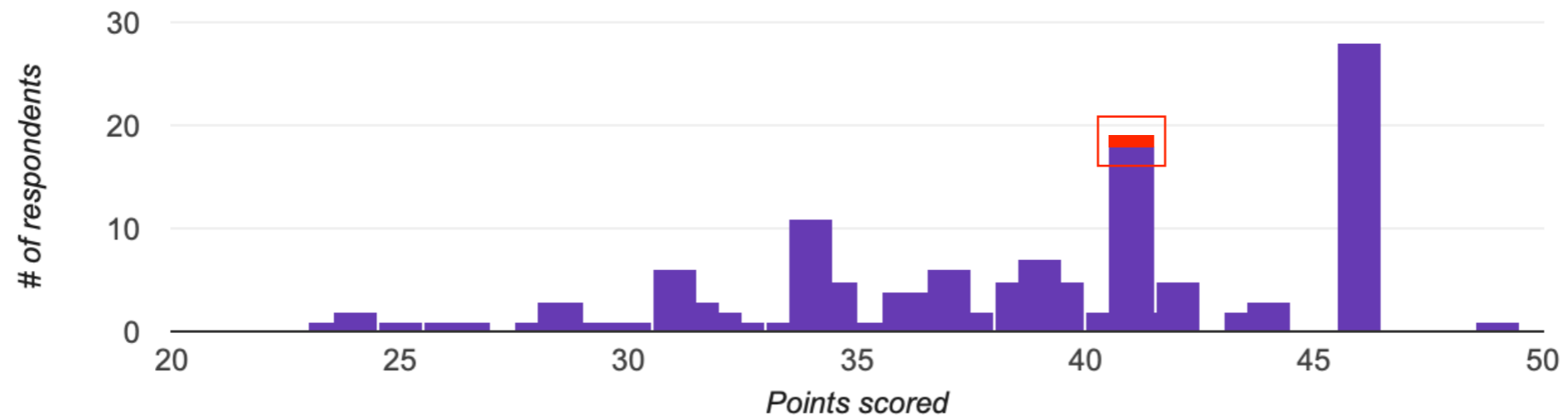


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23.5 - 49 points

**Total points distribution**





Analysis  $\mathcal{M}$  satisfies differential privacy if...

For all  $D_1$  and  $D_2$  which **differ in one individual's data...**

Answer **A** and answer **B** are **indistinguishable**

